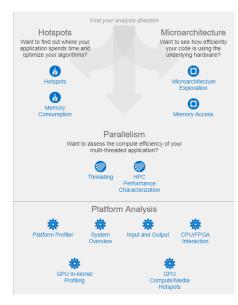
# PRODUCT BRIEF

Performance Profiling Intel® VTune™ Profiler



# **Optimize Software Performance for Modern Hardware**

# Quickly Discover Performance Bottlenecks with Intel® VTune™ Profiler



**Figure 1.** Collect and analyze a rich set of performance data.

### **Collect a Wide Range of Performance Data**

Whether you're tuning a simple application for the first time or doing advanced performance optimization on a threaded MPI application, get the data you need with Intel® VTune™ Profiler. Collect a rich set of performance data for hotspots, threading, locks and waits, DirectX\*, OpenMP\*, Threading Building Blocks, bandwidth, cache, memory access, non-uniform memory, storage latency, OpenCL™ applications, and more.

Profile C, C++, C#, Fortran\*, Python\*, Go\*, Java\* and OpenCL\*—or any mix—running on Intel<sup>®</sup> processors. Unlike single-language profilers, Intel VTune Profiler analyzes mixed code.

- See more data. CPU, FPU, GPU, threading, memory access, and more.
- Get fast answers. Easy analysis turns data into insight.
- Create faster code. Tune with accurate data and low overhead.
- **Choose your workflow.** Select from local/remote collection or command line/ graphical interface.

Analysis Configuration	Collection Log Summary							
Grouping: Function / Call S	Stack						• 🛠 🔉 🖁	
		CPU Time 🔻						
Function / Call Stack	<b></b>	>>	Spin Time			>>	Microarchitecture	
		Effective Time by Utilization Poor Ok Ideal Over		Lock Contention	Other	Overhead Time	Usage	
grid_intersect	3.490s		0s	0s	0s	0s	37.7%	
sphere_intersect	3.004s		0s	0s	0s	0s	41.6%	
GdipDrawImagePointRe	0.431s 📕		0s	0s	0s	0s	100.0%	
grid_bounds_intersect	0.176s 📕		0s	0s	0s	0s	27.9%	
func@0x6b101e50	0.130s 📕		0s	0s	0s	0s	10.9%	
kmp_fork_barrier	Os		0.057s	0s	0.064s	0.001s	100.0%	
▶ < _kmp_launch_threa	Os		0.055s	0s 0.064s		0.001s	100.0%	
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<b>- +</b> : Q	🕊 🖆 6200ms 6300n	ns	6400ms	6500	ms	Thread	•	
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OMP Worker Thread #2 (TI								
OMR Worker Thread #1 (TL)						n and Overhead Ti		
OMP Master Thread #0 (TL.,						cktick Sample		
CPU 1						CPU Tir	ne	
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**Figure 2.** Hotspot analysis shows where your application's time is spent. It can also give a detailed analysis of threading performance showing the potential performance gain and common causes of poor performance such as imbalance, lock contention, forking, scheduling, and reduction.

### Powerful Data Analysis Saves You Time

Good data isn't enough. You need to mine the data for insight. Save time with high-level summaries and powerful analysis to sort, filter, and visualize results on the timeline and on your source.

#### **New for 2020**

- Memory Analysis. Design and optimize for Intel<sup>®</sup> Optane<sup>™</sup> DC persistent memory.
- Platform Profiler. Get a faster, better display of system metrics.
- I/O Analysis. PCIe\* device metrics
- **HPC Analysis.** Vectorization metrics, process and thread affinity, and Lustre\* parallel file I/O metrics (preview feature)
- Snapshot. Communication pattern diagnosis and Open MPI\*
- GPU and OpenCL<sup>™</sup> Application Analysis. Inline filer and instruction count
- Linux\*. More analysis types enabled without adding drivers.
- **Performance Analysis Cookbooks.** User-friendly recipes.
- Improved user interface. New user? Take an overlay tour.
- More descriptive name (formerly Intel<sup>®</sup> VTune<sup>™</sup> Amplifier).

#### Flexible Options: Get It Standalone or Integrated in Comprehensive Development Tool Suites

Intel VTune Profiler is available as a standalone product and as part of other comprehensive software development tool suites.

- Intel® Parallel Studio XE is a development suite to help developers build high-performance, scalable, reliable parallel code for enterprise to cloud, and HPC to AI applications. The suite also includes compilers, libraries, and other analysis tools.
- Intel<sup>®</sup> System Studio Professional and Ultimate editions are used for system and IoT development of smart, connected devices. The suite also includes a compiler, libraries, analysis tools, and cloud connectors and provides access to 400+ sensors.

Get more analysis tools to complement Intel VTune Profiler:

• Intel® Advisor helps optimize vectorization, threading, and flow graphs. Roofline analysis finds loops with the most headroom for improvement.

 Intel<sup>®</sup> Trace Analyzer and Collector examines MPI applications and tells Intel VTune Profiler which loops will benefit most from threading optimization.

#### **Get the Data You Need**

- Hotspot (statistical call tree)
- Thread profiling with locks and waits analysis
- Memory access, cache miss, bandwidth, NUMA analysis
- FLOPS and FPU utilization
- **Storage accesses** mapped to source, latency histogram, I/O wait
- OpenCL program kernel tracing and GPU offload

#### **Easy to Use**

- No special compiles: C, C++, C#, Fortran, Java, Python, Go, ASM\*
- Microsoft Visual Studio\* IDE integration
- Graphical interface and command line
- Local and remote data collection, multi-rank setup for MPI applications
- Collect on Linux, Windows\*, FreeBSD\*, Android\* and select embedded OSs.
- Analyze results on Linux, Windows, and macOS\* hosts.

#### **Find Answers Fast**

- View results on the source/assembly.
- **OpenMP** scalability analysis and graphical frame analysis.
- **Memory analysis:** Tune data structures and optimize NUMA latency.
- Storage analysis: Find I/O bottlenecks.
- Filter out extraneous data with the timeline and viewpoints.
- Visualize thread and task activity on the timeline.

#### Low-Overhead/High-Resolution Hardware Profiling

In addition to basic analysis that works on both Intel and compatible processors, Intel VTune Profiler has advanced analysis that uses the on-chip Performance Monitoring Unit (PMU) on Intel processors to collect data with very low overhead. This also finds important performance issues like cache misses, branch mispredictions, bandwidth, and more.

# **Product Details**

Hotspots analysis gives you a sorted list of the functions using a lot of CPU time. This is where tuning gives you the biggest benefit. Click [+] for the call stacks. Double-click to see the source.

Quickly Locate Code Taking a Lot of CPU Time

Grouping: Function / Call Stack								
	CPU Time <							
Function / Call Stack	Effective Time by Utilization V 🔊 Idle 🔋 Poor 🔋 Ok 🛢 Ideal 🏮 Over	Spin Overhead Time Time						
FireObject::checkCollision	3.348s <b>200</b>	Os Os						
▶ ĸ [Loop at line 1453 in FireObject::P	2.771s 2.771s	Os Os						
▶ < [Loop at line 1491 in FireObject::P	0.578s 📔	0s 0s						
[Loop at line 1453 in FireObject::Proces	1.052s	0s 0s						
▶ rand	0.696s 📒	0s 0s						
▶ ParticleEmitter::FirePatch::initParticle	0.520s 📒	Os Os						

## **Product Details (Continued)**

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							Over Time	Creation	Scheduling	Reduction
and now OpenCL kernels. See line-level								Creation	Scheddling	Reduction
profiling details on the source.		for( u32 j = rang	geBegin;	j < range			0.0%		0.0%	
pronting details on the source.	1,457 { 1,458 FireObject *pfo = m pFireObj		0.0%		0.0%		0.0%			
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	1,460	{			0.0%		0.0%		0.0%	
	1,461	// if it	passes	this test	0.0%		0.0%	0.0%	0.0%	0.09
Tuning OpenMP and Threading Building		· (	10		(E					
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blocks is Easier with the hight but						CF	PU Time 🔻	Spin Time		^
See the cause of threading inefficiencies					tive Time by Uti		Lock		Other	
sorted by potential impact for accurate data		Idle		Idle F	oor 📒 Ok 📳			Contention		Julei
and low overhead.	▶ he	▶ heart_demo (rank 15) 99.64					21.705s	0.009s	8.027s	6.780s
Optimize Multi-Rank Hybrid MPI/OpenMP		▼ heart_demo (rank 17)					21.650s	0.017s	8.012s	6.864s
		[Serial - outside any region]					15.719s	0.004s	6.602s	4.992s
Sort results by impact of improved OpenMP	make_rk_step\$omp\$parallel			21.290s			1.418s	0.003s		0.445s
performance.		make_rk_step\$omp\$parallel: 21.183s					1.431s	0.004s		0.419s
		make_rk_step\$omp\$	parallel:	21.239s			1.320s	0.003s	6 0.350s	0.461s
Identify Which Memory Objects Are										
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Bottlenecks		· ·		-	y Latency			P M		
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Host Operating Systems	Windows, Linux, macOS					
Target Operating Systems	Linux, Windows, FreeBSD, Android, and select embedded operating systems					
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